

# Differential Cryptanalysis of Keccak Variants

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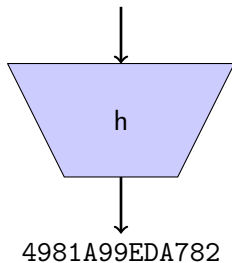
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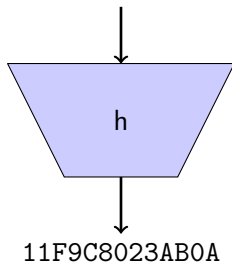
# Cryptographic Hash Functions

“Today is the 18th of December...”



# Cryptographic Hash Functions

“Today is the 19th of December...”



# Cryptographic Hash Functions

## Applications:

- ▶ Message Integrity
- ▶ Digital Signature Schemes
- ▶ Password Protection
- ▶ Key Derivation
- ▶ Payment Schemes (Bitcoin)
- ▶ ...

## Requirements:

- ▶ no secret parameter
- ▶ fast to compute
- ▶ secure

# Cryptographic Hash Functions

## Security Requirements

- ▶ Preimage Resistance:  
Given  $h(x)$  find  $x$
- ▶ Second-Preimage Resistance:  
Given  $x, h(x)$  find  $y \neq x$  s.t.  $h(x) = h(y)$
- ▶ Collision Resistance:  
Find  $x, y$  with  $x \neq y$  s.t.  $h(x) = h(y)$

## Generic Attack

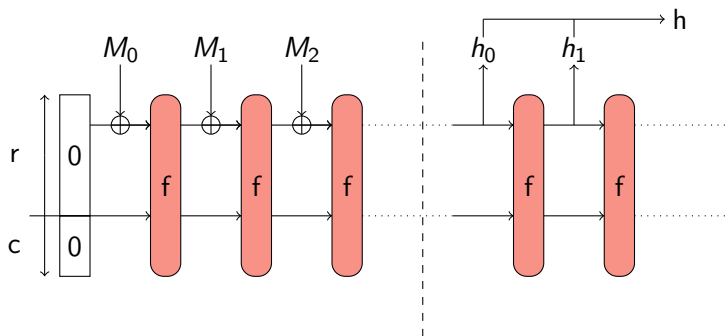
Complexity  $2^n$  for (second) preimage and  $2^{n/2}$  for collisions.

# Keccak

- ▶ Designed by Bertoni, Daemen, Peeter and Van Assche
- ▶ Selected by NIST in October 2012 to become the new SHA-3 standard.
- ▶ Based on the sponge construction.
- ▶ Uses fixed size permutation Keccak-f.
- ▶ Uses 1600-bit permutation for SHA-3.
- ▶ Supports output sizes of  $\{224, 256, 384, 512\}$ -bit.

# Sponge Construction

Takes arbitrary sized input and produces arbitrary sized output.



- ▶ The permutation is of size  $b = r + c$ .
- ▶ Security claim of  $2^{c/2}$

# Sponge Construction

Comparison of Keccak with  $c = 2n$  and  $c = n$ .

	Keccak-256		Keccak-512	
Capacity	512	256	1024	512
Rate	1088	1344	576	1088
Coll. Res.	$2^{128}$	$2^{128}$	$2^{256}$	$2^{256}$
Preimg Res.	$2^{256}$	$2^{128}$	$2^{512}$	$2^{256}$
Performance		+23.5%		+88.9%

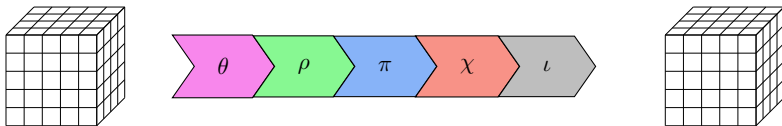


# Keccak

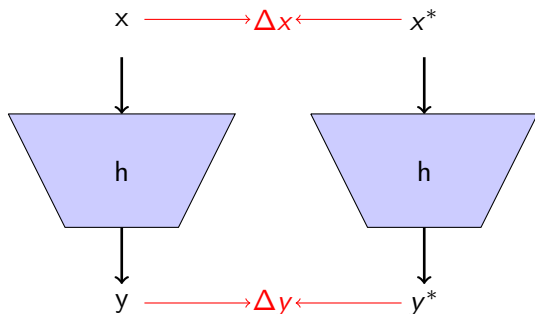
The Keccak-f function

- ▶ 24 rounds
- ▶ Each round is composed of five steps  $\theta, \rho, \pi, \chi, \iota$ .
- ▶ Only XOR, AND, NOT and data-independent rotations are used.

One round of Keccak-f:



# Differential Cryptanalysis



- ▶  $\Delta x \neq 0$  and  $\Delta y = 0$  gives a collision.
- ▶ Find a differential characteristic leading to zero output difference.
- ▶ Find a confirming message pair.

## Related Work

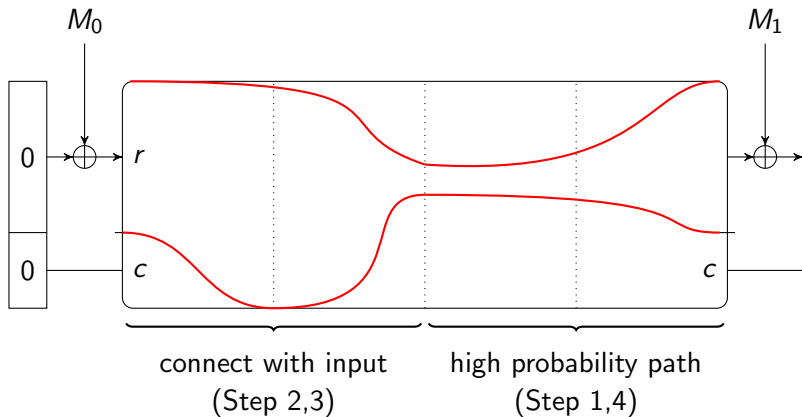
Attack by Naya-Plasencia et al.

- ▶ A 2-round practical attack using high probability paths [NPRM11].

Attack by Dinur et al.

- ▶ A 4-round practical attack on Keccak-224/256 by using the same high probability path [DDS12].
- ▶ Theoretical attacks on 5-round Keccak-256, 4-round Keccak-384 and 3-round Keccak-512 [DDS13].
- ▶ Connect to the starting point using an algebraic method.

# Attack Strategy



# Attack Strategy

## Finding the high probability paths

- ▶ Using linearized model of Keccak
- ▶ Gives a linear code over  $\mathbb{F}_2$
- ▶ Probability that characteristic holds related to the Hamming weight
- ▶ Find codewords with low Hamming weight<sup>1</sup>

Gives us high probability paths leading to (internal) collisions for different Keccak variants.

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<sup>1</sup><http://www.iaik.tugraz.at/content/research/krypto/codingtool/>

## Connecting the paths

Using an automatic search tool to connect the path to the start.

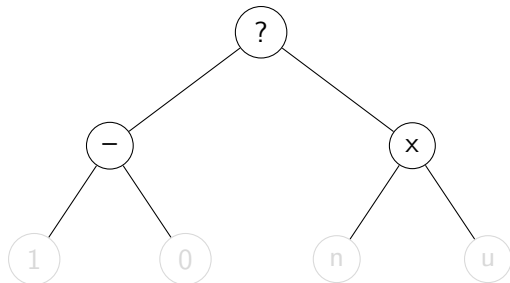
- ▶ Used for instance on SHA-2 [MNS11][MNS13].
- ▶ Guess and determine strategy.

$(X_i, X'_i)$	(0, 0)	(1, 0)	(0, 1)	(1, 1)
?	✓	✓	✓	✓
-	✓			✓
x		✓	✓	
0	✓			
u		✓		
n			✓	
1				✓
⋮				

# Connecting the paths

Search Algorithm [DR06][MNS11]

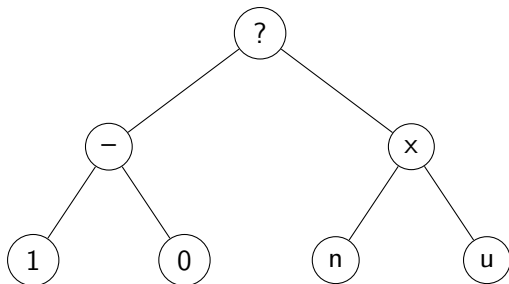
1. **Decision:** Select bit to guess.
2. **Deduction:** Propagate conditions [EMN<sup>+</sup>13].
3. **Backtracking:** Resolve contradictions.



# Connecting the paths

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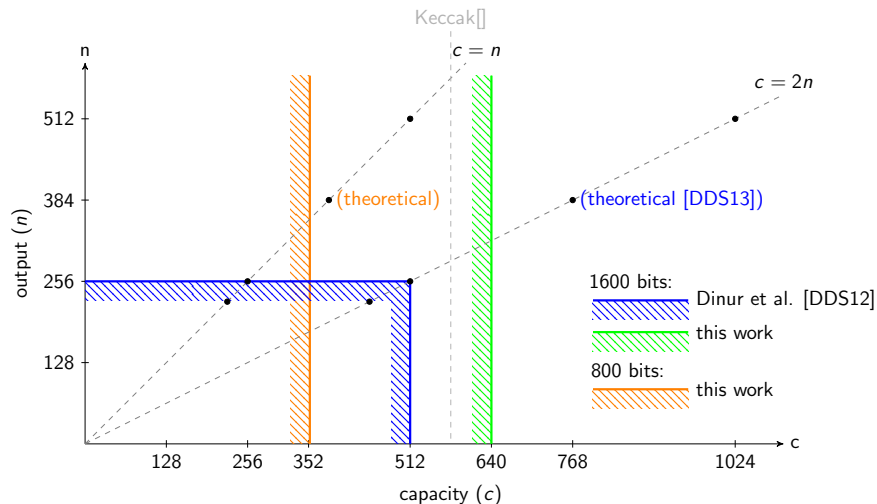
# Example

## State

737bc39f15b62ce3	4-ae-67d9-f67961	72c17e19ecf12b7b	2ba7b749c7949634	fc-cfc935859fb2e
3d196398efcd8-85	fce83de1dec57822	585c3e88-e91a216	7abfed54f57e1dd9	d9a96ed7944d8ede
147b6be6e6-24fdb	--4a7743-1159181	-1df19ab97369543	77a1e8bca7-c--6f	-5e697e1852d7fd5
1a9b2c7d9b5a9abf	2913f4ef6ca6b829	4--b84511fabc4ff	236c8edaa59db4a3	fa16a175b84e4326
6c34feb1242754fb	cb2ea33a4c-db176	b2c5aa5a8-df6238	7bafafd7ee121941	8b4cf1f55781e-9f
96--3182f1fad467	22--9-644fa7e-f-	de--54fb5f2e9a6b	7e--726f824-bd4c	d2--114a6bb11583
96-171-2f1fad467	26--9-644fa7e-f-	de--54fb5f2e9a6b	7e--726f8244b14c	d2--114a6fb51583
96-17112f1fad467	22--b-244fa7e-f-	de--54fb5f2e9a4b	7e--726f8244b14c	d2--114a6bb11583
96-171-2f1fad467	26--9-644fa7e-f2	de--54fb5f2e9a6b	7e--726f884-b14c	d2--114a6bb11583
96-171-2f1fad467	22--9-644fa7e-f-	da--5-fb5f2e9a6b	fe--726f8244b14c	d2--114a6ab11583
---4-8-----	-----4-----	---1-----	-----	-----
---4-8-----	-----4-----	-----	-----	---8-----
-----	-----	-----	-----	-----
-----	-----	---1-----	-----	---8-----
-----	-----	-----	-----	-----
---4-8-----	-----4-----	-----	-----	-----
-----	-----8-----	-----	-----	-----
-----8-----	-----	-----	8-----	-----
-----	-----4-8-----	-----	8-----	-----
---4-----	-----	-----	-----	-----
---4-8-----	---8-----	-----1-----	---8---1-----	---8-4-----
-----	-1---4-----	---4-----	81-----	-----
-----1-8-----	-----	---1---1-----	-----1-----	-----1-----
-----	-----	-----	-----	-----
-----	-----	-----	-----	-----

# Overview

## 4-round attacks on Keccak



# Conclusion

## Results:

- ▶ 4-round practical attack on different Keccak variants.
- ▶ New method to connect paths to the starting point.
- ▶ High probability paths for new variants of Keccak
- ▶ Internal collisions for these variants

Thank you for your attention!

# References I

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